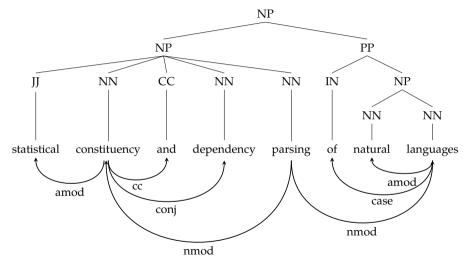
### Statistical Natural Language Processing Statistical Parsing

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University of Tübingen Seminar für Sprachwissenschaft

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### This lecture is about



# Why do we need syntactic parsing?

• Syntactic analysis is an intermediate step in (semantic) interpretation of sentences



As result, it is useful for applications like *question answering, information extraction, ...* 

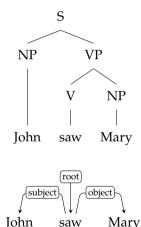
- (Statistical) parsers can also be used as *language models* for applications like *speech recognition* and *machine translation*
- It can be used for *grammar checking*, and can be a useful tool for linguistic research

## Ingredients of a parser

- A grammar
- An algorithm for parsing
- A method for ambiguity resolution

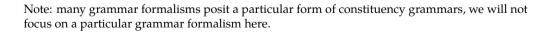
### Dependency vs. constituency

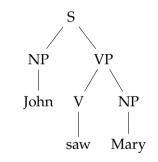
- Constituency grammars are based on units formed by a group of lexical items (constituents or phrases)
- Dependency grammars model binary head-dependent relations between words
- Most of the theory of parsing is developed with constituency grammars
- Dependency grammars has recently become popular in CL



## Constituency grammars

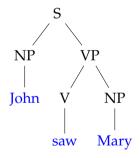
- Constituency grammars are probably the most studied grammars both in linguistics, and computer science
- The main idea is that groups of words form natural groups, or 'constituents', like *noun phrases* or *word phrases*
- *phrase structure grammars* or *context-free grammars* are often used as synonyms





#### A phrase structure grammar is a tuple $(\Sigma, N, S, R)$

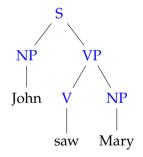
 $\Sigma$  is a set of terminal symbols



A phrase structure grammar is a tuple  $(\Sigma, N, S, R)$ 

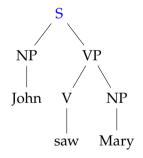
 $\Sigma$  is a set of terminal symbols

N is a set of non-terminal symbols

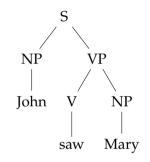


A phrase structure grammar is a tuple  $(\Sigma, N, S, R)$ 

- $\Sigma$  is a set of terminal symbols
- N is a set of non-terminal symbols
- $S \in N$  is a distinguished *start* symbol



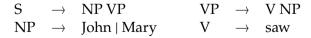
A phrase structure grammar is a tuple ( $\Sigma$ , N, S, R)  $\Sigma$  is a set of terminal symbols N is a set of non-terminal symbols S  $\in$  N is a distinguished *start* symbol R is a set of 'rewrite' rules of the form  $\alpha A\beta \rightarrow \gamma$  for  $A \in N$   $\alpha, \beta, \gamma \in \Sigma \cup N$ 

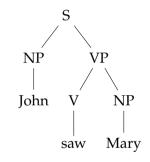


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A phrase structure grammar is a tuple  $(\Sigma,\,N,\,S,\,R)$ 

- $\boldsymbol{\Sigma}~$  is a set of terminal symbols
- $N \ \mbox{is a set of non-terminal symbols}$
- $S \in N$  is a distinguished *start* symbol
- $\begin{array}{ll} R \ \ \text{is a set of 'rewrite' rules of the form} \\ \alpha A\beta \rightarrow \gamma \quad \text{for } A\in N \quad \alpha,\beta,\gamma\in \Sigma\cup N \end{array}$
- The grammar accepts a sentence if it can be derived from S with the rewrite rules R





### Example derivation

The example grammar:

- Phrase structure grammars derive a sentence with successive application of rewrite rules.
   S ⇒NP VP ⇒John VP ⇒John V NP ⇒John saw NP ⇒John saw Mary or, S ⇒John saw Mary
- The intermediate forms that contain non-terminals are called sentential forms

### Constituency grammars and parsing

- Context-free grammars are parseable in  $\mathsf{O}(n^3)$  time complexity using dynamic programming algorithms
- Mildly context-sensitive grammars can also be parsed in polynomial time  $\left(O(n^6)\right)$
- Polynomial time algorithms are not always fast enough in practice
  - We often use approximate solutions with greedy search algorithms

### Where do grammars come from

- Grammars for (statistical) parsing can be either
  - hand crafted (many years of expert effort)
  - extracted from *treebanks* (which also require lots of effort)
  - 'induced' from raw data (interesting, but not as successful)
- Current practice relies mostly on treebanks
- Hybrid approaches also exist
- Grammar induction is not common (for practical models) but exploiting unlabled data is also a common trend

# Context free grammars

recap

- Context free grammars are sufficient for expressing most phenomena in natural language syntax
- Most of the parsing theory (and practice) is build on parsing CF languages
- The context-free rules have the form

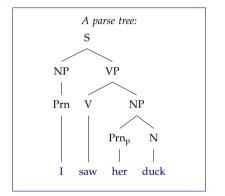
 $A \to \alpha$ 

where A is a single non-terminal symbol and  $\alpha$  is a (possibly empty) sequence of terminal or non-terminal symbols

### An example context-free grammar

$S \rightarrow NP VP$	Derivation of sentence 'she s	aw a duo	ck′		
$\begin{array}{llllllllllllllllllllllllllllllllllll$	$\begin{array}{ll} S & \Rightarrow NP VP \\ NP \Rightarrow Prn \\ Prn \Rightarrow she \\ VP \Rightarrow V NP \\ V \Rightarrow saw \\ NP \Rightarrow Det N \\ Det \Rightarrow a \\ N & \Rightarrow duck \end{array}$	NP   Prn   she	S V saw	VP N Det a	JP N   duck
$Det \rightarrow a \mid the$					

### Representations of a context-free parse tree



A history of derivations:

- $S \Rightarrow NP VP$
- NP  $\Rightarrow$ Prn
- $\Pr n \Rightarrow I$
- VP  $\Rightarrow$ V NP
- $V \Rightarrow saw$
- NP  $\Rightarrow$  Prn<sub>p</sub> N
- $Prn_p \Rightarrow her$
- N  $\Rightarrow$ duck

A sequence with (labeled) brackets  
$$\left[ {}_{S} \left[ {}_{NP} \left[ {}_{Prn} I \right] \right] \left[ {}_{VP} \left[ {}_{V} saw \right] \left[ {}_{NP} \left[ {}_{Prn_{p}} her \right] \left[ {}_{N} duck \right] \right] \right] \right]$$

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### Parsing as search

- Parsing can be seen as search constrained by the grammar and the input
- Top down: start from S, find the derivations that lead to the sentence
- Bottom up: start from the sentence, find series of derivations (in reverse) that leads to S
- Search can be depth first or breadth first for both cases

### Problems with search procedures

- Top-down search considers productions incompatible with the input, and cannot handle left recursion
- Bottom-up search considers non-terminals that would never lead to S
- Repeated work because of backtracking
- ightarrow The result is exponential time complexity in the length of the sentence

Some of these problems can be solved using *dynamic programming*.

## CKY algorithm

- The CKY (Cocke–Younger–Kasami), or CYK, parsing algorithm is a dynamic programming algorithm
- It processes the input *bottom up*, and saves the intermediate results on a *chart*
- Time complexity for *recognition* is  $O(n^3)$  (with a space complexity of  $O(n^2)$ )
- It requires the CFG to be in Chomsky normal form (CNF)

### Chomsky normal form (CNF)

- A CFG is in CNF, if the rewrite rules are in one of the following forms
  - $\ A \to B \ C$
  - $\ A \to a$

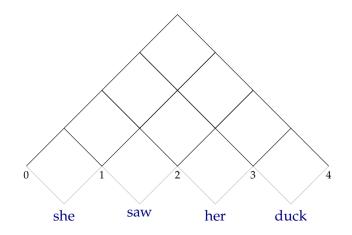
where A, B, C are non-terminals and a is a terminal

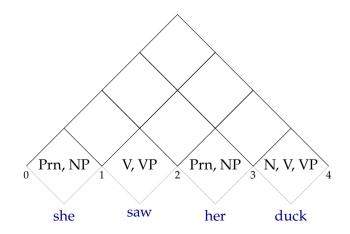
- Any CFG can be converted to CNF
- Resulting grammar is *weakly equivalent* to the original grammar:
  - it generates/accepts the same language
  - but the derivations are different

### Converting to CNF: example

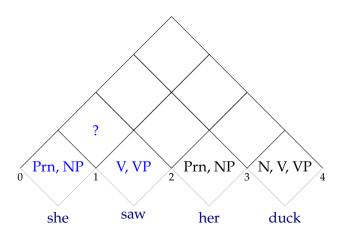
- For rules with > 2 RHS symbols S  $\rightarrow$  Aux NP VP  $\Rightarrow$  S  $\rightarrow$  Aux X X  $\rightarrow$  NP VP
- For rules with < 2 RHS symbols NP  $\rightarrow$ Prn  $\Rightarrow$  NP  $\rightarrow$  she | her

 $S \rightarrow NP VP$  $S \rightarrow Aux NP VP$  $NP \rightarrow Det N$  $NP \rightarrow Prn$  $NP \rightarrow NP PP$  $VP \rightarrow V NP$  $VP \ \rightarrow V$  $VP \rightarrow VP PP$  $PP \rightarrow Prp NP$  $N \rightarrow duck$  $N \rightarrow park$  $N \rightarrow parks$  $V \rightarrow duck$  $V \rightarrow ducks$ V  $\rightarrow$  saw  $Prn \rightarrow she \mid her$  $Prp \rightarrow in | with$ Det  $\rightarrow$  a | the



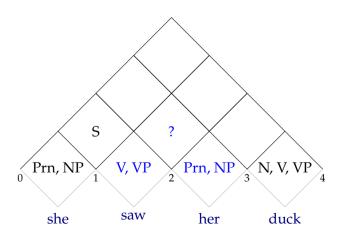


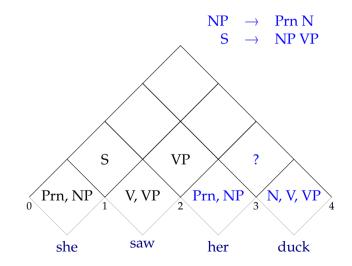


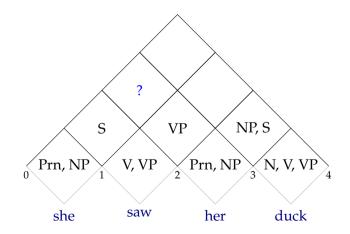


recognition example

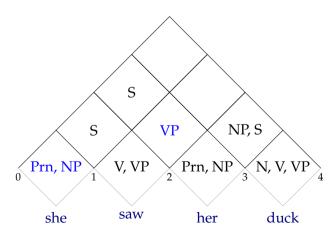
 $VP \quad \rightarrow \quad V \ NP$ 

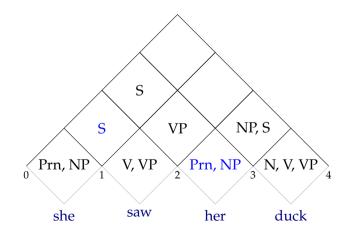


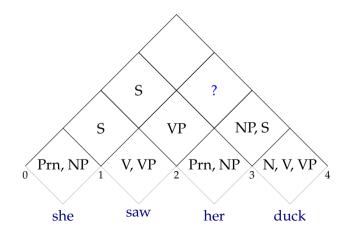


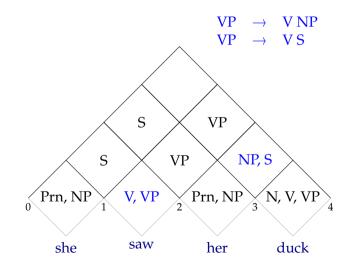


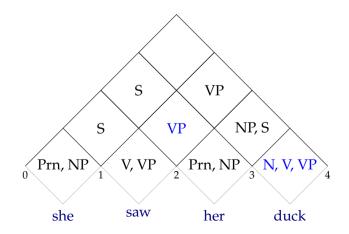


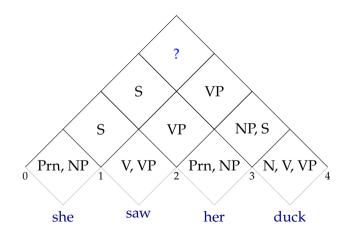






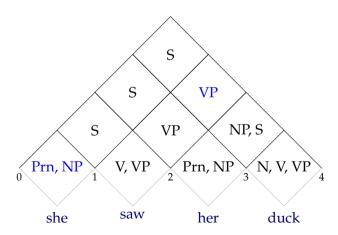


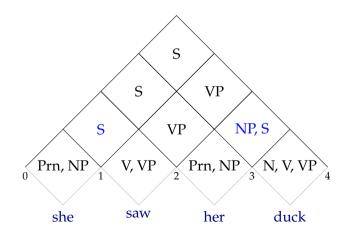




recognition example

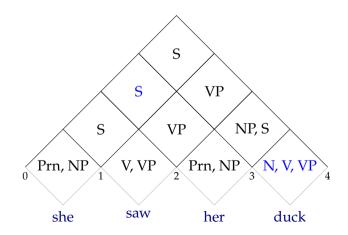
 $S \quad \to \quad NP \; VP$ 





### **CKY** demonstration

recognition example



### CKY demonstration: the chart

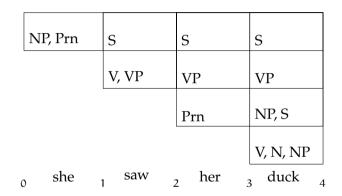
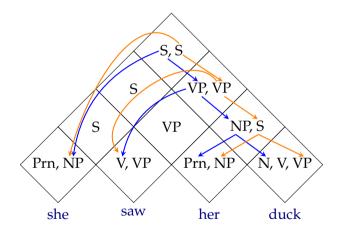


Chart is a 2-dimensional array, hence  $O(n^2)$  space complexity.

### Parsing requires back pointers



## CKY summary

- + CKY avoids re-computing the analyses by storing the earlier analyses (of sub-spans) in a table
- It still computes lower level constituents that are not allowd by the grammar
- CKY requires the grammar to be in CNF
- CKY has O(n<sup>3</sup>) recognition complexity
- For parsing we need to keep track of backlinks
- CKY can effciently store all possible parses in a chart
- Enumerating all possible parses have exponential complexity (worst case)

## Earley algorithm

- Earley algorithm is a top down parsing algorithm
- It allows arbitrary CFGs
- Keeps record of constituents that are

predicted using the grammar (top-down) in-progress with partial evidence completed based on input seen so far at every position in the input string

• Time complexity is  $O(n^3)$ 

## Summary: context-free parsing algorithms

- Naive search for parsing is intractable
- Dynamic programming algorithms allow polynomial time recognition
- Parsing may still be exponential in the worse case
- Ambiguity: CKY or Earley parse tables can represent ambiguity, but cannot say anything about which parse is the best

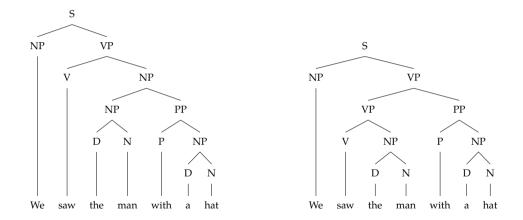
## Pretty little girl's school (again)



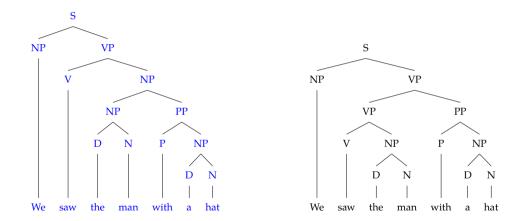
Cartoon Theories of Linguistics, SpecGram Vol CLIII, No 4, 2008. http://specgram.com/CLIII.4/school.gif

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#### The task: choosing the most plausible parse



#### The task: choosing the most plausible parse



## Statistical parsing

- Find the most plausible parse of an input string given all possible parses
- We need a scoring function, for each parse, given the input
- We typically use probabilities for scoring, task becomes finding the parse (or tree), t, given the input string *w*

$$t_{\text{best}} = \arg\max_{t} P(t \mid \boldsymbol{w})$$

• Note that some ambiguities need a larger context than the sentence to be resolved correctly

## Probabilistic context free grammars (PCFG)

A probabilistic context free grammar is specified by,

- $\Sigma$  is a set of terminal symbols
- $N \ \mbox{is a set of non-terminal symbols}$
- $S \in N$  is a distinguished *start* symbol
- $\mathbb R\;$  is a set of rules of the form

$$A \rightarrow \alpha$$
 [p]

where A is a non-terminal,  $\alpha$  is string of terminals and non-terminals, and p *is the probability associated with the rule* 

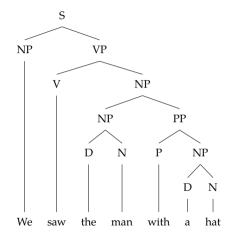
- The grammar accepts a sentence if it can be derived from S with rules  $R_1 \ldots R_k$
- The probability of a parse t of input string  $\bm{w}, P(t\,|\,\bm{w}),$  corresponding to the derivation  $R_1\ldots R_k$  is

$$P(t \,|\, \boldsymbol{w}) = \prod_{1}^{k} p(R_{i})$$

where  $p(R_{\rm i})$  is the probability of the rule  $R_{\rm i}$ 

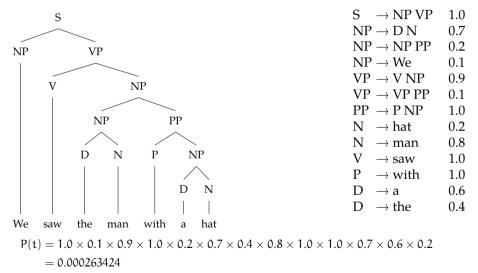
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# PCFG example (1)



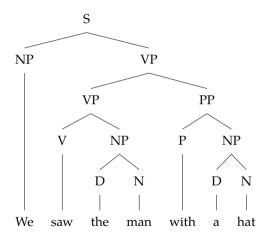
S	$\rightarrow$ NP VP	1.0
-		2.10
INP	$P \rightarrow D N$	0.7
NP	$P \to \mathrm{NP} \ \mathrm{PP}$	0.2
NP	$^{\prime}  ightarrow \mathrm{We}$	0.1
VP	ightarrow V NP	0.9
VP	ightarrow VP PP	0.1
PP	ightarrow P NP	1.0
Ν	ightarrow hat	0.2
Ν	$\rightarrow$ man	0.8
V	$\rightarrow$ saw	1.0
Р	ightarrow with	1.0
D	$\rightarrow$ a	0.6
D	$\rightarrow$ the	0.4

# PCFG example (1)



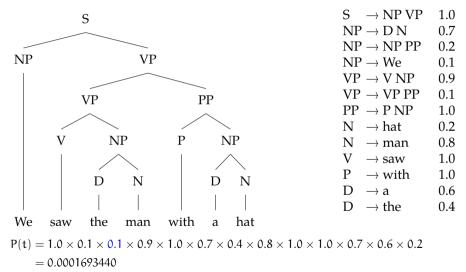
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## PCFG example (2)



S	ightarrow NP VP	1.0
NP	$^{\prime} \rightarrow \mathrm{D}\mathrm{N}$	0.7
NP	$P \to \mathrm{NP} \ \mathrm{PP}$	0.2
NP	$^{\prime}  ightarrow \mathrm{We}$	0.1
VP	$\rightarrow V \ NP$	0.9
VP	$\rightarrow \mathrm{VP} \ \mathrm{PP}$	0.1
$\mathbf{PP}$	ightarrow P NP	1.0
Ν	ightarrow hat	0.2
Ν	$\rightarrow$ man	0.8
V	$\rightarrow$ saw	1.0
Р	ightarrow with	1.0
D	$\rightarrow$ a	0.6
D	$\rightarrow$ the	0.4

## PCFG example (2)



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### Where do the rule probabilities come from?

- Supervised: estimate from a treebank, e.g., using maximum likelihood estimation
- Unsupervised: expectation-maximization (EM)

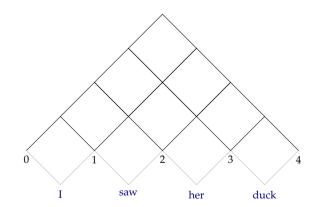
#### PCFGs - an interim summary

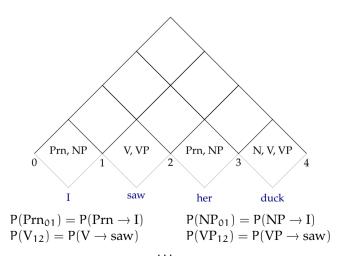
- PCFGs assign probabilities to parses based on CFG rules used during the parse
- PCFGs assume that the rules are independent
- PCFGs are generative models, they assign probabilities to P(t, *w*), we can calcuate the probability of a sentence by

$$P(\boldsymbol{w}) = \sum_{t} P(t, \boldsymbol{w}) = \sum_{t} P(t)$$

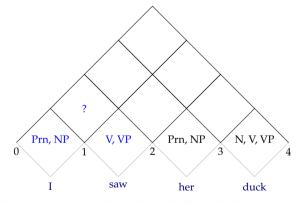
## PCFG chart parsing

- Both CKY and Earley algorithms can be adapted to PCFG parsing
- CKY matches PCFG parsing quite well
  - to get the best parse, store the constituent with the highest probability in every cell of the chart
  - to get n-best best parse (beam search), store the n-best constituents in every cell in the chart



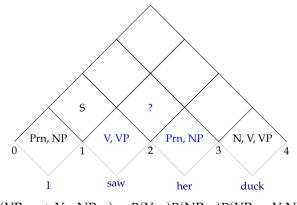




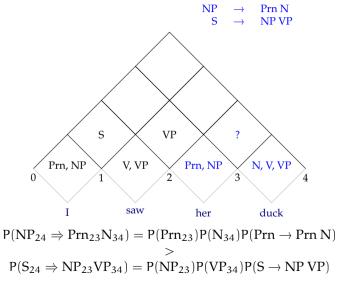


 $\mathsf{P}(\mathsf{S}_{02} \Rightarrow \mathsf{NP}_{01}\mathsf{VP}_{12}) = \mathsf{P}(\mathsf{NP}_{01})\mathsf{P}(\mathsf{VP}_{12})\mathsf{P}(\mathsf{S} \rightarrow \mathsf{NP}\;\mathsf{VP})$ 



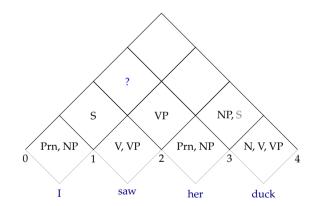


 $\mathsf{P}(\mathsf{VP}_{13} \Rightarrow \mathsf{V}_{12}\mathsf{NP}_{23}) = \mathsf{P}(\mathsf{V}_{12})\mathsf{P}(\mathsf{NP}_{23})\mathsf{P}(\mathsf{VP} \to \mathsf{V} \ \mathsf{NP})$ 

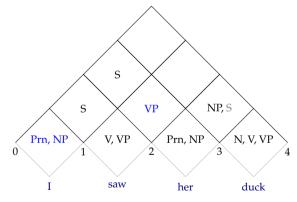


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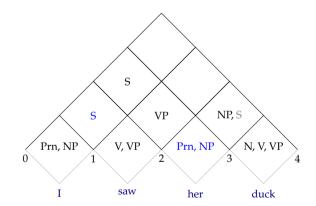
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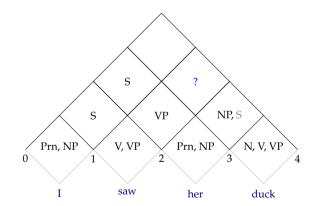


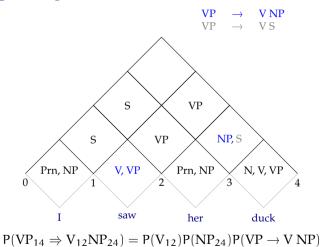


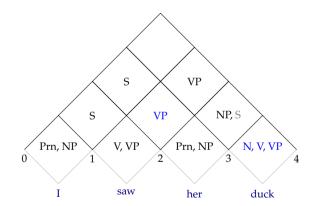


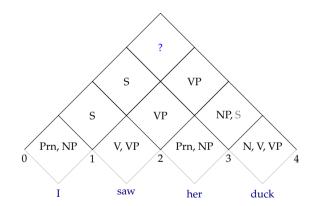
 $\mathsf{P}(\mathsf{S}_{03} \Rightarrow \mathsf{NP}_{01}\mathsf{VP}_{23}) = \mathsf{P}(\mathsf{NP}_{01})\mathsf{P}(\mathsf{VP}_{13})\mathsf{P}(\mathsf{S} \rightarrow \mathsf{NP}\,\mathsf{VP})$ 



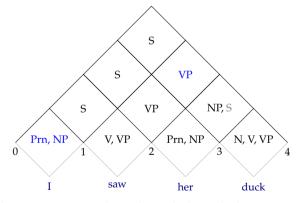




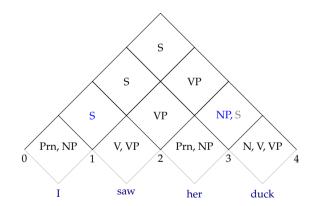


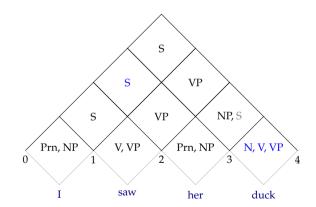






 $\mathsf{P}(S_{14} \Rightarrow NP_{01}VP_{14}) = \mathsf{P}(NP_{01})\mathsf{P}(VP_{14})\mathsf{P}(S \rightarrow NP~VP)$ 





#### What makes the difference in PCFG probabilities?

$S \Rightarrow NP VP$	1.0	$S \Rightarrow NP VP$	1.0
$NP \Rightarrow We$	0.1	$NP \Rightarrow We$	0.1
$VP \Rightarrow VP PP$	0.1	$VP \Rightarrow V NP$	0.7
$\mathrm{VP} \Rightarrow \mathrm{V}\mathrm{NP}$	0.8	$V \Rightarrow saw$	1.0
$V \Rightarrow saw$	1.0	$NP \Rightarrow NP PP$	0.2
$\mathrm{NP} \Rightarrow \mathrm{D}\mathrm{N}$	0.7	$\mathrm{NP} \Rightarrow \mathrm{D} \ \mathrm{N}$	0.7
$D \Rightarrow the$	0.4	$D \Rightarrow the$	0.4
$N \Rightarrow man$	0.8	$N \Rightarrow man$	0.8
$PP \Rightarrow P NP$	1.0	$PP \Rightarrow P NP$	1.0
$P \Rightarrow with$	1.0	$P \Rightarrow with$	1.0
$\text{NP} \Rightarrow \text{D} \text{ N}$	0.7	$\text{NP} \Rightarrow \text{D} \text{ N}$	0.7
$D \Rightarrow a$	0.6	$D \Rightarrow a$	0.6
$N \Rightarrow hat$	0.2	$N \Rightarrow hat$	0.2

#### What makes the difference in PCFG probabilities?

	1.0		1.0
$S \Rightarrow NP VP$	1.0	$S \Rightarrow NP VP$	1.0
$\text{NP} \Rightarrow \text{We}$	0.1	$NP \Rightarrow We$	0.1
$VP \Rightarrow VP PP$	0.1	$\mathrm{VP} \Rightarrow \mathrm{V} \ \mathrm{NP}$	0.7
$\mathrm{VP} \Rightarrow \mathrm{V}\mathrm{NP}$	0.8	$V \Rightarrow saw$	1.0
$V \Rightarrow saw$	1.0	$NP \Rightarrow NP PP$	0.2
$NP \Rightarrow D \ N$	0.7	$\text{NP} \Rightarrow \text{D} \text{ N}$	0.7
$D \Rightarrow the$	0.4	$D \Rightarrow the$	0.4
$N \Rightarrow man$	0.8	$N \Rightarrow man$	0.8
$PP \Rightarrow P NP$	1.0	$PP \Rightarrow P NP$	1.0
$P \Rightarrow with$	1.0	$P \Rightarrow with$	1.0
$NP \Rightarrow D \ N$	0.7	$\text{NP} \Rightarrow \text{D} \text{ N}$	0.7
$D \Rightarrow a$	0.6	$D \Rightarrow a$	0.6
$N \Rightarrow hat$	0.2	$N \Rightarrow hat$	0.2

The parser's choice would not be affected by lexical items!

#### What is wrong with PCFGs?

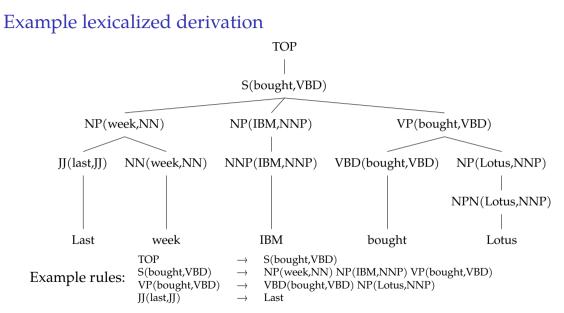
- In general: the assumption of independence
- The parents affect the correct choice for children, for example, in English  $NP \to Prn$  is more likely in the subject position
- The lexical units affect the correct decision, for example:
  - We eat the pizza with hands
  - We eat the pizza with mushrooms
- Additionally: PCFGs use local context, difficult to incorporate arbitrary/global features for disambiguation

### Solutions to PCFG problems

- Independence assumptions can be relaxed by either
  - Parent annotation
  - Lexicalization
- To condition on arbitrary/global information: discriminative models
- Most practical PCFG parsers are lexicalized, and often use a re-ranker conditioning on other (global) features

# Lexicalizing PCFGs

- Replace non-terminal X with X(h), where h is a tuple with the lexical word and its POS tag
- Now the grammar can capture (head-driven) lexical dependencies
- But number of nonterminals grow by  $|V|\times |T|$
- Estimation becomes difficult (many rules, data sparsity)
- Some treebanks (e.g., Penn Treebank) do not annotate heads, they are automatically annotated (based on heuristics)



### Evaluating the parser output

- A parser can be evaluated
  - $extrinsically\ based on its effect on a task (e.g., machine translation) where it is used$

intrinsically based on the match with ideal parsing

- The typically evaluation (intrinsic) is based on a gold standard (GS)
- Exact match is often
  - very difficult to achieve (think about a 50-word newspaper sentence)
  - not strictly necessary (recovering parts of the parse can be useful for many purposes)

#### Parser evaluation metrics

• Common evaluation metrics are (PARSEVAL):

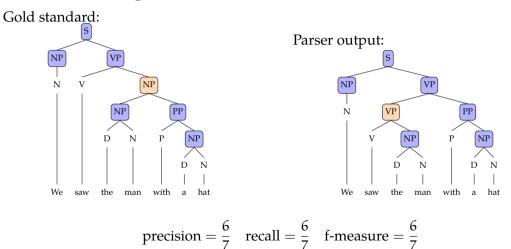
precision the ratio of correctly predicted nodes
 recall the nodes (in GS) that are predicted correctly
f-measure harmonic mean of precision and recall (2×precision×recall
precision+recall)

• The measures can be

unlabled the spans of the nodes are expected to match labeled the node label should also match

- Crossing brackets (or average non-crossing brackets)
  - ( We ( saw ( them ( with binoculars )))) ( We (( saw them ) ( with binoculars )))
- Measures can be averaged per constituent (micro average), or over sentences (macro average)

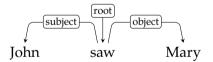
## PARSEVAL example



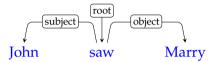
#### Problems with PARSEVAL metrics

- PARSEVAL metrics favor certain type of structures
  - Results are surprisingly well for flat tree structures (e.g., Penn treebank)
  - Results of some mistakes are catastrophic (e.g., low attachment)
- Not all mistakes are equally important for semantic distinctions
- Some alternatives:
  - Extrinsic evaluation
  - Evaluation based on extracted dependencies

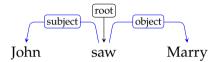
- Dependency grammars gained popularity in (particularly in computational) linguistics rather recently, but their roots can be traced back to a few thousand years
- The main idea is capturing the relation between the words, rather than grouping them into (abstract) constituents



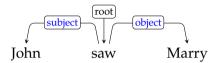
Note: like constituency grammars, we will not focus on a particular dependency formalism, but discuss it in general in relation to parsing.



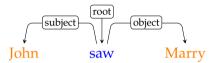
• No constituents, units of syntactic structure are words



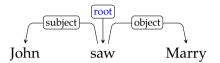
- No constituents, units of syntactic structure are words
- The structure of the sentence is represented by asymmetric binary relations between syntactic units



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- The links (relations) have labels (dependency types)



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- The links (relations) have labels (dependency types)
- Each relation defines one of the words as the head and the other as dependent



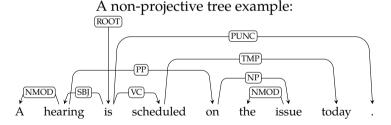
- No constituents, units of syntactic structure are words
- The structure of the sentence is represented by asymmetric binary relations between syntactic units
- The links (relations) have labels (dependency types)
- Each relation defines one of the words as the head and the other as dependent
- Often an artificial root node is used for computational convenience

#### Projective vs. non-projective dependencies

- If a dependency graph has no crossing edges, it is said to be *projective*, otherwise *non-projective*
- Non-projectivity stems from long-distance dependencies and free word order

#### Projective vs. non-projective dependencies

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# Parsing with dependency grammars

- Projective parsing can be done in polynomial time
- Non-projective parsing is NP-hard (without restrictions)
- For both, it is a common practice to use greedy (e.g., linear time) algorithms

## Dependency grammar: definition

A dependency grammar is a tuple (V, A)

- V is a set of nodes corresponding to the (syntactic) words (we implicitly assume that words have indexes)
- A is a set of arcs of the form  $(w_i, r, w_j)$  where
  - $w_i \in V$  is the head
    - r is the type of the relation (arc label)
  - $w_j \in V$  is the dependent

This defines a directed graph.

#### Dependency grammars: common assumptions

- Every word has a single head
- The dependency graphs are acyclic
- The graph is connected
- With these assumptions, the representation is a tree
- Note that these assumptions are not universal but common for dependency parsing

# Dependency parsing

- Dependency parsing has many similarities with context-free parsing (e.g., trees)
- They also have some different properties (e.g., number of edges and depth of trees are limited)
- Dependency parsing can be
  - grammar-driven (hand crafted rules or constraints)
  - data-driven (rules/model is learned from a treebank)
- There are two main approaches:

Graph-based similar to context-free parsing, search for the best tree structure Transition-based similar to shift-reduce parsing (used for programming language parsing), but using greedy search for the best transition sequence

### Transition based parsing

- Inspired by shift-reduce parsing, single pass over the input
- Use a stack and a buffer of unprocessed words
- Parsing as predicting a sequence of transitions like
  - LEFT-ARC: mark current word as the head of the word on top of the stack RIGHT-ARC: mark current word as a dependent of the word on top of the stack SHIFT: push the current word to the stack
- Algorithm terminates when all words in the input are processed
- The transitions are not naturally deterministic, best transition is predicted using a machine learning method

(Yamada and Matsumoto 2003; Nivre, Hall, and Nilsson 2004)

## A typical transition system



 $\text{Left-Arc}_{r}: \ (\sigma|w_{\mathfrak{i}},w_{\mathfrak{j}}|\beta,A) \Rightarrow (\sigma \quad ,w_{\mathfrak{j}}|\beta,A \cup \{(w_{\mathfrak{j}},r,w_{\mathfrak{i}})\})$ 

- pop *w*<sub>i</sub>
- add arc  $(w_j, r, w_i)$  to A
- keep  $w_i$  in the buffer

 $\text{Right-Arc}_r: \ (\sigma|w_i,w_j|\beta,A) \Rightarrow (\sigma \quad,w_i|\beta,A \cup \{(w_i,r,w_j)\})$ 

- pop  $w_i$
- add arc  $(w_i, r, w_j)$  to A
- move  $w_i$  to the buffer ( $w_j$  is removed from the buffer)

Shift:  $(\sigma , w_j | \beta, A) \Rightarrow (\sigma | w_j, \beta, A)$ 

• push *w*<sub>j</sub> to the stack (remove it from the buffer)

(Kübler, McDonald, and Nivre 2009, p.23)

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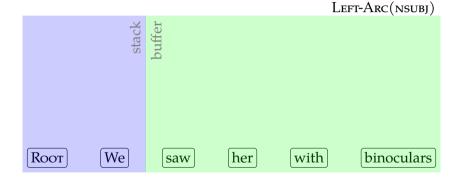
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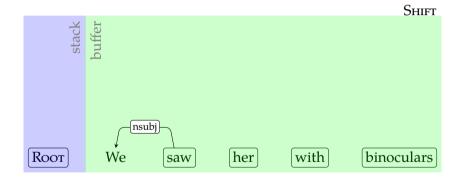
### Transition based parsing: example



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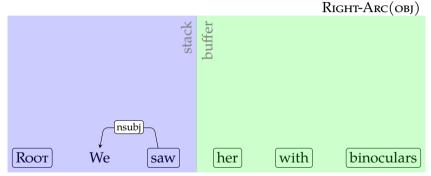
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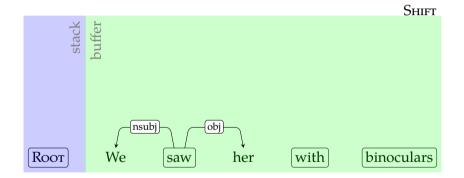
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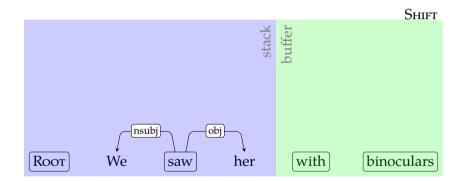
### Transition based parsing: example



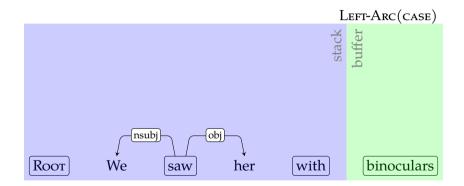
Note: We need Shift for NP attachment.

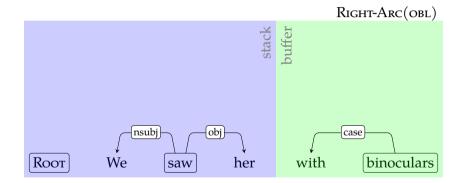
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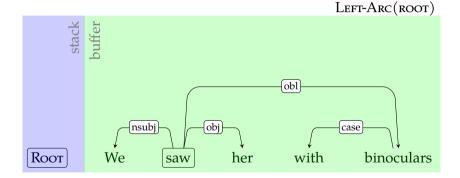




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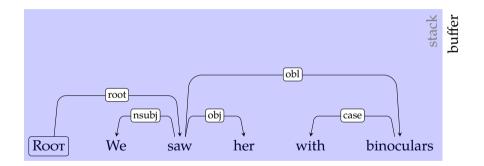




Grammars Constituency grammars & parsing Dependency grammars & parsing Summary

#### Transition based parsing: example

Shift



### Making transition decisions

- In shift-reduce parsing the actions are deterministic
- In transition-based dependency parsing, we need to choose among all possible transitions
- The typical method is to train a (discriminative) classifier on features extracted from gold-standard *transition sequences*
- Almost any machine learning method is applicable. Common choices include
  - Memory-based learning
  - Support vector machines
  - (Deep) neural networks

### Features for transition-based parsing

- The features come from certain 'addresses' in the parser configuration, for example
  - The word at the stack top (or nth from stack top)
  - The first/second word on the buffer
  - Right/left dependents of the word on top of the stack/buffer
- For each possible 'address', we can make use of features like
  - Word form, lemma, POS tag, morphological features, word embeddings
  - Dependency relations  $(w_i, r, w_j)$  triples
- Note that for some 'address'-'feature' combinations and in some configurations the values may be missing

# The training data

- We want features like,
  - lemma[Stack] = duck
  - POS[Stack] = NOUN
  - ...
- But treebank gives us:

1	Read	read	VERB	VB	Mood=Imp VerbForm=Fin	0	root
2	on	on	ADV	RB	-	1	advmod
3	to	to	PART	TO	-	4	mark
4	learn	learn	VERB	VB	VerbForm=Inf	1	xcomp
5	the	the	DET	DT	Definite=Def	6	det
6	facts	fact	NOUN	NNS	Number=Plur	4	obj
7		•	PUNCT		-	1	punct

• The treebank has the outcome of the parser, but not the features we need

## The training data

- The features for transition-based parsing have to be from *parser configurations*
- The data (treebanks) need to be preprocessed for obtaining the training data
- Construct a transition sequence by parsing the sentences, and using treebank annotations (the set A) as an 'oracle'
- Decide for

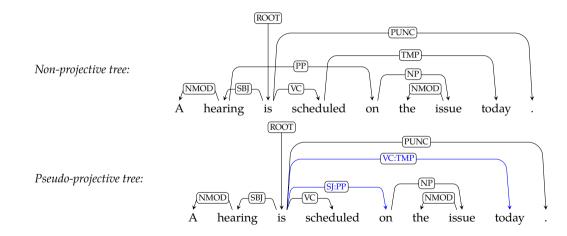
```
Left-Arc<sub>r</sub> if (\beta[0], r, \sigma[0]) \in A
Right-Arc<sub>r</sub> if (\sigma[0], r, \beta[0]) \in A
and all dependents of \beta[0] are attached
Shift otherwise
```

• There may be multiple sequences that yield the same dependency tree, the above defines a 'canonical' transition sequence

## Non-projective parsing

- The transition-based parsing we defined so far works only for projective dependencies
- One way to achieve (limited) non-projective parsing is to add special Left-Arc and Right-Arc transitions to/from non-top words from the stack
- Another method is pseudo-projective parsing:
  - preprocessing to 'projectivize' the trees before training
    - The idea is to attach the dependents to a higher level head that preserves projectivity, while marking it on the new dependency label
  - postprocessing for restoring the projectivity after parsing
    - Re-introduce projectivity for the marked dependencies

## Pseudo-projective parsing



## Transition based parsing: summary/notes

- Linear time, greedy parsing
- Can be extended to non-projective dependencies
- One can use arbitrary features
- We need some extra work for generating gold-standard transition sequences from treebanks
- Early errors propagate, transition-based parsers make more mistakes on long-distance dependencies
- The greedy algorithm can be extended to beam search for better accuracy (still linear time complexity)

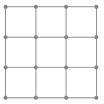
# Graph-based parsing: preliminaries

- Enumerate all possible dependency trees
- Pick the best scoring tree
- Features are based on limited parse history (like CFG parsing)
- Two well-known flavors:
  - Maximum (weight) spanning tree (MST)
  - Chart-parsing based methods

Eisner 1996; McDonald et al. 2005

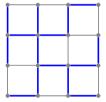
### MST parsing: preliminaries Spanning tree of a graph

• Spanning tree of a connected graph is a sub-graph which is a tree and traverses all the nodes



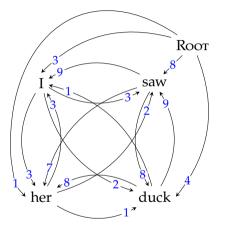
## MST parsing: preliminaries Spanning tree of a graph

- Spanning tree of a connected graph is a sub-graph which is a tree and traverses all the nodes
- For fully-connected graphs, the number of spanning trees are exponential in the size of the graph
- The problem is well studied
- There are efficient algorithms for enumerating and finding the optimum spanning tree on weighted graphs

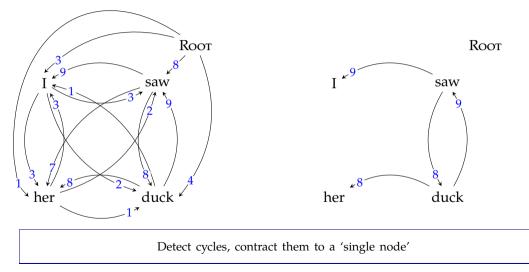


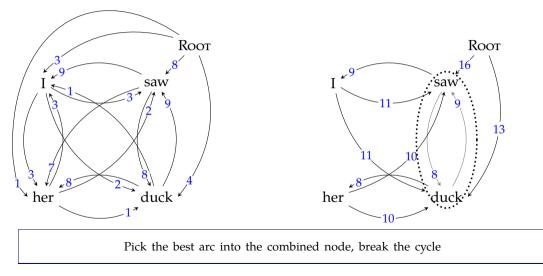
# MST algorithm for dependency parsing

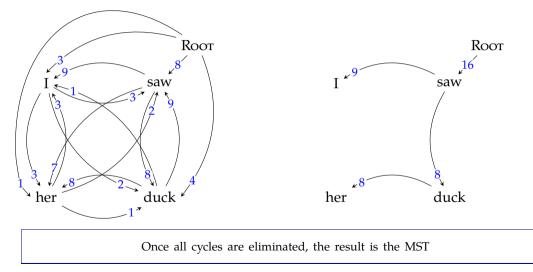
- For directed graphs, there is a polynomial time algorithm that finds the minimum/maximum spanning tree (MST) of a fully connected graph (Chu-Liu-Edmonds algorithm)
- The algorithm starts with a dense/fully connected graph
- Removes edges until the resulting graph is a tree



For each node select the incoming arc with highest weight







# Properties of the MST parser

- The MST parser is non-projective
- There is an alrgorithm with  $O(n^2)$  time complexity  $_{\scriptscriptstyle (Tarjan\,1977)}$
- The time complexity increases with typed dependencies (but still polynomial)
- The weights/parameters are associated with edges (often called 'arc-factored')
- We can learn the arc weights directly from a treebank
- However, it is difficult to incorporate non-local features

# CKY for dependency parsing

- The CKY algorithm can be adapted to projective dependency parsing
- For a naive implementation the complexity increases drastically  $O(n^6)$ 
  - Any of the words within the span can be the head
  - Inner loop has to consider all possible splits
- For projective parsing, the observation that the left and right dependents of a head are independently generated reduces the comlexity to  $O(n^3)$

(Eisner 1997)

#### Non-local features

- The graph-based dependency parsers use edge-based features
- This limits the use of more global features
- Some extensions for using 'more' global features are possible
- This often leads non-projective parsing to become intractable

#### External features

- For both types of parsers, one can obtain features that are based on unsupervised methods such as
  - clustering
  - dense vector representations (embeddings)
  - alignment/transfer from bilingual corpora/treebanks

(Koo, Carreras, and Collins 2008)

## Errors from different parsers

- Different parsers make different errors
  - Transition based parsers do well on local arcs, worse on long-distance arcs
  - Graph based parsers tend to do better on long-distance dependencies
- Parser combination is a good way to combine the powers of different models. Two common methods
  - Majority voting: train parsers separately, use the weighted combination of their results
  - Stacking: use the output of a parser as features for another

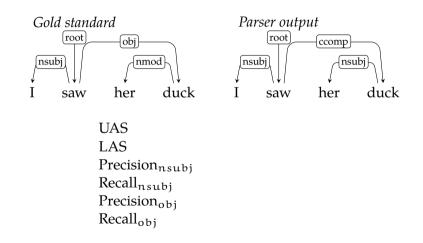
(McDonald and Satta 2007; Sagae and Lavie 2006; Nivre and McDonald 2008)

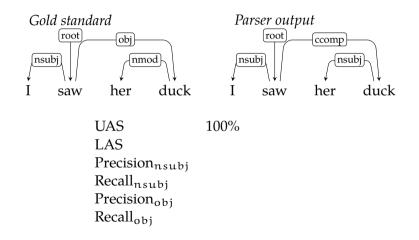
## Evaluation metrics for dependency parsers

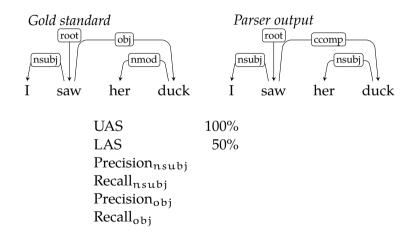
- Like CF parsing, exact match is often too strict
- *Attachment score* is the ratio of words whose heads are identified correctly.
  - Labeled attachment score (LAS) requires the dependency type to match
  - Unlabeled attachment score (UAS) disregards the dependency type
- *Precision/recall/F-measure* often used for quantifying success on identifying a particular dependency type

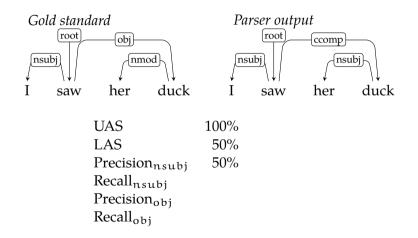
precision is the ratio of correctly identified dependencies (of a certain type) recall is the ratio of dependencies in the gold standard that parser predicted correctly

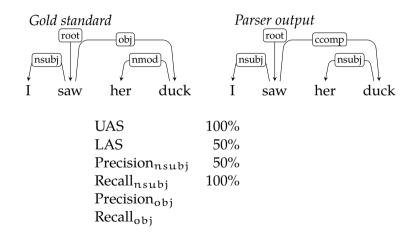
f-measure is the harmonic mean of precision and recall  $\left(\frac{2 \times \text{precision} \times \text{recall}}{\text{precision} + \text{recall}}\right)$ 

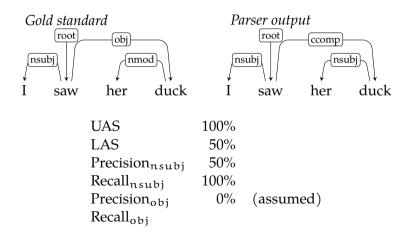


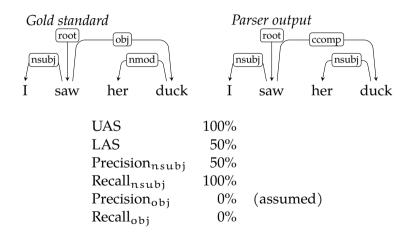












## Averaging evaluation scores

- As in context-free parsing, average scores can be macro-average or sentence-based micro-average or word-based
- Consider a two-sentence test set with

	words	correct
sentence 1	30	10
sentence 2	10	10

- word-based average attachment score:
- sentence-based average attachment score:

## Averaging evaluation scores

- As in context-free parsing, average scores can be macro-average or sentence-based micro-average or word-based
- Consider a two-sentence test set with

	words	correct
sentence 1	30	10
sentence 2	10	10

- word-based average attachment score: 50% (20/40)
- sentence-based average attachment score: 66% ((1 + 1/3)/2)

## Dependency parsing: summary

- Dependency relations are often easier to interpret
- It is also claimed that dependency parsers are more suitable for parsing free-word-order languages
- Dependency relations are between words, no phrases or other abstract nodes are postulated
- Two general methods:

transition based greedy search, non-local features, fast, less accurate graph based exact search, local features, slower, accurate (within model limitations)

- Combination of different methods often result in better performance
- Non-projective parsing is more difficult
- Most of the recent parsing research has focused on better machine learning methods (mainly using neural networks)

#### Next week

Ion/Wed Wrap-up/summary Fri Exam

# Where to go from here?

- Textbook includes good coverage of constituency grammars and parsing, online 3rd edition includes a chapter on dependency parsing as well
- The book by Kübler, McDonald, and Nivre (2009) is an accessible introduction to (statistical) dependency parsing
- For more on linguistic and mathematical foundations of parsing:
  - Müller (2016) is a new open-source text book on Grammar formalisms.
  - Aho and Ullman (1972) is the classical reference (available online) for parsing (programming languages) and also includes discussion of grammar classes in the Chomsky hierarchy. A more up-to-date alternative is Aho, Lam, et al. (2007).
  - There is a brief introductory section on dependency grammars in Kübler, McDonald, and Nivre (2009), for a classical reference see tesniere2015, English translation of the original version (Tesnière 1959).

# Pointers to some treebanks

Treebanks are the main resource for statistical parsing. A few treebank-related resources to have a look at until next time:

• Universal dependencies project, documentation, treebanks:

http://universaldependencies.org/

• Tübingen treebanks:

TüBa-D/Z written German

TüBa-D/S spoken German

TüBa-E/S spoken English

TüBa-J/S spoken Japanese

available from http://www.sfs.uni-tuebingen.de/en/ascl/resources/corpora.html

- TüNDRA a treebank search and visualization application with the above treebanks and few more
  - Main version:

https://weblicht.sfs.uni-tuebingen.de/Tundra/

– New version (beta):

https://weblicht.sfs.uni-tuebingen.de/tundra-beta/

# CKY algorithm

```
\begin{array}{l} \textbf{function CKY}(words, grammar) \\ \textbf{for } j \ \leftarrow \ 1 \ \textbf{to } \ \text{LENGTH}(words) \ \textbf{do} \\ \texttt{table}[j-1,j] \leftarrow \{A|A \rightarrow words[j] \in \texttt{grammar}\} \\ \textbf{for } i \ \leftarrow \ j-1 \ \textbf{downto} \ 0 \ \textbf{do} \\ \textbf{for } k \ \leftarrow \ i+1 \ \textbf{to} \ j-1 \ \textbf{do} \\ \texttt{table}[i,j] \ \leftarrow \ \texttt{table}[i,j] \cup \\ \{A|A \rightarrow BC \in \texttt{grammar} \ \textbf{and} \\ B \in \texttt{table}[i,k] \ \textbf{and} \\ C \in \texttt{table}[k,j]\} \end{array}
```

return table

# Even more examples

(newspaper headlines)

- FARMER BILL DIES IN HOUSE
- TEACHER STRIKES IDLE KIDS
- SQUAD HELPS DOG BITE VICTIM
- BAN ON NUDE DANCING ON GOVERNOR'S DESK
- PROSTITUTES APPEAL TO POPE
- KIDS MAKE NUTRITIOUS SNACKS
- DRUNK GETS NINE MONTHS IN VIOLIN CASE
- MINERS REFUSE TO WORK AFTER DEATH

saw

а

duck

S	ightarrow NP VP
S	$\to Aux \: X$
Х	$\rightarrow$ NP VP
NP	$\rightarrow Det \ N$
NP	ightarrow she   her
NP	$\rightarrow$ NP PP
VP	ightarrow V NP
VP	$\rightarrow$ duck saw
VP	ightarrow VP PP
$\mathbf{PP}$	$ ightarrow \operatorname{Prp}\operatorname{NP}$
Ν	$\rightarrow$ duck
Ν	$\rightarrow$ park
Ν	$\rightarrow$ parks
V	$\rightarrow$ duck
V	ightarrow ducks
V	$\rightarrow$ saw
Prn	ightarrow she   her
Prp	$\rightarrow$ in   with
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she

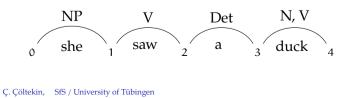
saw <sub>2</sub> a <sub>3</sub> duck <sub>4</sub>

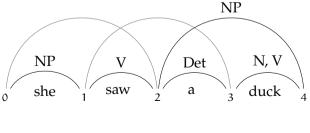
S	$\rightarrow$ NP VP
S	$\to Aux \: X$
Х	ightarrow NP VP
NP	$\rightarrow Det \ N$
NP	ightarrow she   her
NP	$\rightarrow$ NP PP
VP	ightarrow V NP
VP	$\rightarrow$ duck saw
VP	ightarrow VP PP
$\mathbf{PP}$	$ ightarrow \operatorname{Prp}\operatorname{NP}$
Ν	$\rightarrow$ duck
Ν	$\rightarrow$ park
Ν	$\rightarrow$ parks
V	$\rightarrow$ duck
V	ightarrow ducks
V	$\rightarrow$ saw
Prn	$\rightarrow$ she   her
Prp	ightarrow in   with
Det	-Surfamer thenester 2020

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she 1

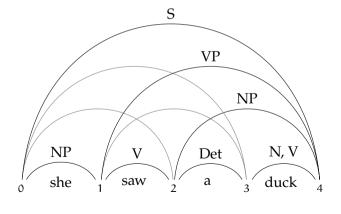
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S  $\rightarrow$  NP VP S  $\rightarrow$  Aux X Х  $\rightarrow$  NP VP  $NP \rightarrow Det N$  $NP \rightarrow she \mid her$  $NP \rightarrow NP PP$  $VP \rightarrow V NP$  $VP \rightarrow duck |saw|...$  $VP \rightarrow VP PP$  $PP \rightarrow Prp NP$  $N \rightarrow duck$  $N \rightarrow park$  $N \rightarrow parks$  $V \rightarrow duck$ V  $\rightarrow$  ducks V  $\rightarrow$  saw  $Prn \rightarrow she \mid her$  $Prp \rightarrow in \mid with$ Det -suramethenester 2020



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- S  $\rightarrow$  NP VP S  $\rightarrow Aux X$  $X \rightarrow NP VP$  $NP \rightarrow Det N$  $NP \rightarrow she \mid her$  $NP \rightarrow NP PP$  $VP \rightarrow V NP$  $VP \rightarrow duck |saw|...$  $VP \rightarrow VP PP$  $PP \rightarrow Prp NP$  $N \rightarrow duck$  $N \rightarrow park$  $N \rightarrow parks$  $V \rightarrow duck$  $V \rightarrow ducks$ V  $\rightarrow$  saw  $Prn \rightarrow she \mid her$  $Prp \rightarrow in \mid with$
- Det -Summerhemester 2020 A.5